# Difference Geometry 

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Unit Geometry and CAD
University Innsbruck
July 22-23, 2010


Lecture 1:
Introduction

## Three disciplines

## Differential geometry

- infinitesimally neighboring objects
- calculus, applied to geometry


## Difference geometry

- finitely separated objects
- elementary geometry instead of calculus

Discrete differential geometry

- "modern" difference geometry
- emphasis on similarity and analogy to differential geometry



## History

1920-1970 H. Graf, R. Sauer, W. Wunderlich:

- didactic motivation
- emphasis on flexibility questions
since 1995 U. Pinkall, A. I. Bobenko and many others:
- deep theory (arguably richer than the smooth case)
- development of organizing principles (Bobenko and Suris, 2008)
- connections to integrable systems
- applications in physics, computer graphics, architecture, ...


## Motivation for a discrete theory

Didactic reasons:

- easily accessible and concrete
- requires little a priori knowledge (advanced calculus vs. elementary geometry)
Rich theory: - at least as rich as smooth theory
- clear explanations for "mysterious" phenomena in the smooth setting
Applications: - high potential for applications due to discretizations
- numerous open research questions


## Overview

Lecture 1: Introduction
Lecture 2: Discrete curves and torses
Lecture 3: Discrete surfaces and line congruences
Lecture 4: Discrete curvature lines
Lecture 5: Parallel nets, offset nets and curvature
Lecture 6: Cyclidic net parametrization

## Literature

A. I. Bobenko, Yu. B. Suris

Discrete Differential Geometry. Integrable Structure American Mathematical Society (2008)
$\otimes$ R. Sauer
Differenzengeometrie
Springer (1970)
Further references to literature will be given during the lecture and posted on the web-page
http://geometrie.uibk.ac.at/schroecker/difference-geometry/

## Software

Adobe Reader Recent versions that can handle 3D-data. http://get.adobe.com/jp/reader
Rhinoceros 3D-CAD; evaluation version (fully functional, save limit) is available at http://rhino3d.com.
Geogebra Dynamic 2D geometry, open source. Download at http://geogebra.org.
Cabri 3D Dynamic 3D geometry. Evaluation version (restricted mode after 30 days) available at http://cabri.com/cabri-3d.html.
Maple Symbolic and numeric calculations.
Worksheets will be made available in alternative formats. http://maplesoft.com
Asymptote Graphics programming language used for most pictures in this lecture. http://asymptote.sourceforge.net

## Conventions for this lecture

- If not explicitly stated otherwise, we assume generic position of all geometric entities.
- Concepts from differential geometry are used as motivation. Results are usually given without proof.
- Concepts from elementary geometry are usually visualized and named. You can easily find the proofs on the internet.
- Concepts from other fields (projective geometry, CAGD etc.) will be explained in more detail upon request.
- Questions are highly appreciated.


## An example from planar kinematics

One-parameter motion

$$
\alpha: I \subset \mathbb{R} \rightarrow \mathrm{SE}(2), \quad t \mapsto \alpha(t)=\alpha_{t}
$$

where

$$
\alpha_{t}: \Sigma \rightarrow \Sigma^{\prime}, \quad x \mapsto \alpha_{t}(x)=x(t)
$$

and

$$
\alpha_{t}(x)=\left(\begin{array}{rr}
\cos \varphi(t) & -\sin \varphi(t) \\
\sin \varphi(t) & \cos \varphi(t)
\end{array}\right) \cdot\binom{x_{1}}{x_{2}}+\binom{a_{1}(t)}{a_{2}(t)}
$$

## The cycloid (circle rolls on line)



$$
\begin{gathered}
\varphi(t)=-t, \quad a_{1}(t)=t, \quad a_{2}(t)=0 \\
\alpha_{t}(x)=\left(\begin{array}{rr}
\cos (-t) & -\sin (-t) \\
\sin (-t) & \cos (-t)
\end{array}\right) \cdot\binom{x_{1}}{x_{2}}+\binom{t}{0}
\end{gathered}
$$

## Corresponding result from three positions theory

Theorem (Inflection circle)
The locus of points $x$ such that the trajectory $x(t)=\alpha_{t}(x)$ has an inflection point at $t=t_{0}$ is a circle.

## - inflection-circle.mw

## Theorem

Given are three positions $\Sigma_{0}, \Sigma_{1}$, and $\Sigma_{2}$ of a moving frame $\Sigma$ in the Euclidean plane $\mathbb{R}^{2}$. Generically, the locus of points $x \in \Sigma$ such that the three corresponding points $x_{0} \in \Sigma_{0}, x_{1} \in \Sigma_{1}, x_{2} \in \Sigma_{2}$ are collinear is a circle.

## Corresponding result from three positions theory



The line $y_{1} \vee y_{2} \vee y_{3}$ is the Simpson line to $x$.

## Comparison

## Smooth theorem

- Formulation requires knowledge (planar kinematics, inflection point, ...)
- Proof requires calculus and algebra (differentiation, circle equation)

Discrete theorem

- elementary formulation and proof
- smooth theorem by limit argument


## The cycloid evolute



Theorem
The locus of curvature centers of the cycloid (its evolute) is a congruent cycloid
$\rightarrow$ cycloid.3dm

## The discrete cycloid evolute



## Theorem (see Hoffmann 2009)

n even: The locus of circle centers through three consecutive points of a discrete cycloid (its vertex evolute) is a congruent discrete cycloid.
$\boldsymbol{n}$ odd: The locus of circle centers tangent to three consecutive edges of a discrete cycloid (its edge evolute) is a congruent discrete cycloid.

## Literature

Inflection circle: Chapter 8, §9 of Bottema and Roth (1990).
Discrete cycloid: Hoffmann (2009)
Simpson line: Bottema (2008).
$\theta$
O. Bottema

Topics in Elementary Geometry
Springer (2008)
$\otimes$ O. Bottema, B. Roth
Theoretical Kinematics
Dover Publications (1990)
$\otimes$ T. Hoffmann
Discrete Differential Geometry of Curves and Surfaces
Faculty of Mathematics, Kyushu University (2009)

## Lecture 2: <br> Discrete Curves and Torses

## Smooth and discrete curves

## Smooth curve:

$$
\gamma: I \subset \mathbb{R} \rightarrow \mathbb{R}^{d}, \quad u \mapsto \gamma(u)
$$

Regularity condition:

$$
\frac{\mathrm{d} \gamma}{\mathrm{~d} u}(u)=\dot{\gamma}(u) \neq 0
$$

Discrete curve:
$\gamma: I \subset \mathbb{Z} \rightarrow \mathbb{R}^{d}, \quad i \mapsto \gamma(i)=: \gamma_{i}$,
Regularity condition:

$$
\delta \gamma_{i}:=\gamma_{i+1}-\gamma_{i} \neq 0
$$

## Smooth and discrete curves

## Smooth curve:

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Regularity condition:

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Regularity condition:

$$
\delta \gamma_{i}:=\gamma_{i+1}-\gamma_{i} \neq 0
$$

Shift notation:

$$
\begin{gathered}
\gamma_{i} \approx \gamma, \quad \gamma_{i+1} \approx \gamma_{1}, \quad \gamma_{i-1} \approx \gamma_{-1} \\
\quad \text { for example } \delta \gamma=\gamma_{1}-\gamma
\end{gathered}
$$

## Example

Discuss the regularity of

$$
\gamma(t)=\binom{t-\sin t}{1-\cos t}
$$

Solution

$$
\dot{\gamma}(t)=\binom{1-\cos t}{\sin t}=0 \Longleftrightarrow t=2 k \pi, k \in \mathbb{Z}
$$

## Example

Derive a parametrization of the discrete cycloid and discuss its regularity.

Solution

$$
\begin{gathered}
\gamma_{k}=\sum_{l=0}^{k}\left(1-e^{-\mathrm{i} l \frac{2 \pi}{n}}\right)=\sum_{l=0}^{k}\left(\binom{1}{0}-\binom{\cos \frac{2 l \pi}{n}}{\sin \frac{2 l \pi}{n}}\right) \\
\gamma_{k}-\gamma_{k-1}=1-e^{-\mathrm{i} k \frac{2 \pi}{n}}=0 \Longleftrightarrow \frac{k}{n} \in \mathbb{Z}
\end{gathered}
$$

## Tangent, principal normal, and bi-normal


tangent vector $t:=\delta \gamma /\|\delta \gamma\|$
normal vector

- $n \perp t$,
- $\|n\|=1$,
- $n$ parallel to $\gamma_{-1} \vee \gamma \vee \gamma_{1}$,
- same orientation of $t_{-1} \times t$ and $t \times n$


## Tangent, principal normal, and bi-normal



Definition
The Frenet-frame is the orthonormal frame with origin $\gamma$ and axis vectors $t, n, b$.

## Tangent, principal normal, and bi-normal



Osculating plane: incident with $\gamma$, orthogonal to $b$ Normal plane: incident with $\gamma$, orthogonal to $t$ Rectifying plane: incident with $\gamma$, orthogonal to $n$

## Smooth and discrete curvature

## Smooth curvature:

Infinitesimal change of tangent direction with respect to arc length:

$$
\varkappa(t)=\frac{\|\dot{\gamma}(t) \times \ddot{\gamma}(t)\|}{\|\dot{\gamma}(t)\|^{3}}
$$

Discrete curvature:

$$
\varkappa:=\frac{\sin \varphi}{s} \quad \text { where } \quad \varphi=\varangle\left(t_{-1}, t\right), s=\left\|\gamma_{1}-\gamma\right\|
$$

- Assume $\varphi \in\left[0, \frac{\pi}{2}\right]$ or $\varphi \in\left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$ (in case of $d=2$ ).
- We will later encounter different notions of curvature.


## Smooth and discrete torsion

Smooth torsion: Change of bi-normal direction with respect to arc length (measure of "planarity"):

$$
\tau(t)=\frac{\langle\dot{\gamma}(t) \times \ddot{\gamma}(t), \dddot{\gamma}(t)\rangle}{\|\dot{\gamma}(t) \times \ddot{\gamma}(t)\|^{2}}
$$

Discrete torsion:

$$
\tau:=\frac{\sin \psi}{s} \quad \text { where } \quad \psi=\varangle\left(b, b_{1}\right)
$$

- assume $\psi \in\left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$
- $\psi \geqslant 0 \Longleftrightarrow$ helical displacement of Frenet frame at $\gamma_{-1}$ to Frenet frame at $\gamma$ is a right screw
- $\tau \equiv 0 \Longleftrightarrow$ curve is planar


## Infinite sequence of refinements

- Assume that all points $\gamma$ are sampled from a smooth curve $\gamma(s)$, parametrized by arc-length.
- Consider an infinite sequence of refinements $\gamma_{i}=\gamma(\varepsilon i), \varepsilon \rightarrow 0$.

Curvature: $\varkappa \rightarrow \varkappa(s)$
Torsion: $\tau \rightarrow \tau(s)$
Frenet frame: $t \rightarrow t(s)=\frac{\dot{\gamma}(s)}{\|\dot{\gamma}(s)\|}, n \rightarrow n(s), b \rightarrow b(s)$

## The fundamental theorems of curve theory

## Theorem

Curvature $\varkappa(s)$ and torsion $\tau(s)$ as functions of the arc length determine a space curve up to rigid motion.

Proof.
Existence and uniqueness of an initial value problem for a system of partial differential equations.

Theorem
The three functions

- $\varkappa: \mathbb{Z} \rightarrow\left[0, \frac{\pi}{2}\right]$ (curvature),
- $\tau: \mathbb{Z} \rightarrow\left[-\frac{\pi}{2} \frac{\pi}{2}\right]$ (torsion), and
- $s: \mathbb{Z} \rightarrow \mathbb{R}^{+}$(arc-length)
uniquely determine a discrete space curve up to rigid motion.
Proof.
Elementary construction.


## The fundamental theorems of curve theory

Theorem
Curvature $\varkappa(s)$ and torsion $\tau(s)$ as functions of the arc length determine a space curve up to rigid motion.

## Proof.

Existence and uniqueness of an initial value problem for a system of partial differential equations.

Corollary
The two functions

- $\varkappa: \mathbb{Z} \rightarrow\left[-\frac{\pi}{2}, \frac{\pi}{2}\right]$ (curvature) and
- $s: \mathbb{Z} \rightarrow \mathbb{R}^{+}$(arc-length)
determine a discrete planar curve.


## Discrete Frenet-Serret equations

$$
\begin{gathered}
t-t_{-1}=(1-\cos \varphi) t+\sin \varphi n \Longrightarrow \\
\frac{t-t_{-1}}{s}=\frac{1-\cos \varphi}{s} t+\varkappa n \\
\frac{1-\cos \varphi}{s}=\frac{\sin \varphi}{s} \cdot \tan \frac{\varphi}{2}=\varkappa \tan \frac{\varphi}{2} \rightarrow 0 \\
t^{\prime}(s):=\frac{\mathrm{d} t}{\mathrm{~d} s}(s)=\lim _{\varepsilon \rightarrow 0} \frac{t-t_{-1}}{s}=\lim _{\varepsilon \rightarrow 0} \varkappa n=\varkappa(s) n(s)
\end{gathered}
$$

## Discrete Frenet-Serret equations

$$
\begin{gathered}
b_{1}-b=(\cos \psi-1) b-\sin \psi n \Longrightarrow \\
\frac{b_{1}-b}{s}=\frac{\cos \psi-1}{s}-\tau n
\end{gathered}
$$

$$
\begin{gathered}
\frac{\cos \psi-1}{s}=-\frac{\sin \psi}{s} \cdot \tan \frac{\psi}{2}=-\tau \cdot \tan \frac{\psi}{2} \rightarrow 0 \\
b^{\prime}(s):=\frac{\mathrm{d} b}{\mathrm{~d} s}(s)=\lim _{\varepsilon \rightarrow 0} \frac{b_{1}-b}{s}=\lim _{\varepsilon \rightarrow 0}-\tau n=-\tau(s) n(s) .
\end{gathered}
$$

## Smooth Frenet-Serret equations

$$
\begin{aligned}
t^{\prime}(s)=\varkappa(s) n(s), \quad b^{\prime}(s)=-\tau(s) n(s) \Longrightarrow \\
\begin{aligned}
n^{\prime}(s):=\frac{\mathrm{d} n}{\mathrm{~d} s}(s) & =-\frac{\mathrm{d}(t \times b)}{\mathrm{d} s}(s) \\
& =-t^{\prime}(s) \times b(s)-t(s) \times b^{\prime}(s) \\
& =-\varkappa(s) n(s) \times b(s)+t(s) \times \tau(s) n(s) \\
& =-\varkappa(s) t(s)+\tau(s) b(s) . \\
\left(\begin{array}{c}
t^{\prime} \\
n^{\prime} \\
b^{\prime}
\end{array}\right) & =\left(\begin{array}{ccc}
0 & \varkappa & 0 \\
-\varkappa & 0 & \tau \\
0 & -\tau & 0
\end{array}\right) \cdot\left(\begin{array}{l}
t \\
n \\
b
\end{array}\right)
\end{aligned}
\end{aligned}
$$

## Discrete torses

## Definition

A discrete torse is a map $T$ from $\mathbb{Z}$ to the space of planes in $\mathbb{R}^{3}$.
, discrete-screw-torse.3dm

rulings: $\ell=T_{-1} \cap T$
edge of regression: $\gamma=T_{-1} \cap T \cap T_{1}$

- $\ell$ is an edge of $\gamma$
- $\ell$ and $\ell_{1}$ intersect
- $T$ is osculating plane of $\gamma$
$\rightsquigarrow$ equivalent definitions based on planes, points, and lines

Application: Design of closed folded strips


## Application: Design of closed folded strips


http://www.archiwaste.org/?p=1109
Institut für Konstruktion und Gestaltung, Universität Innsbruck:
Rupert Maleczek, Eda Schaur
Archiwaste:
Guillaume Bounoure, Chloe Geneveaux

## Literature

R. Sauer's book contains

- the derivation of the Frenet-Serret equations as presented here and
- a treatise on discrete torses.

Q R. Sauer
Differenzengeometrie Springer (1970)

## Lecture 3: <br> Discrete Surfaces and Line Congruences

## Smooth parametrized surfaces

$$
\begin{gathered}
f: U \subset \mathbb{R}^{2} \rightarrow \mathbb{R}^{3}, \quad(u, v) \mapsto f(u, v) \\
f_{u} \times f_{v} \neq 0 \quad \text { where } f_{u}:=\frac{\partial f}{\partial u}, f_{v}:=\frac{\partial f}{\partial v} \\
\text { (tangent vectors to parameter lines) }
\end{gathered}
$$

## Example

Discuss the regularity of the parametrized surface

$$
f(u, v)=\left(\begin{array}{c}
\cos u \cos v \\
\cos u \sin v \\
\sin u
\end{array}\right), \quad(u, v) \in\left(-\frac{\pi}{2}, \frac{\pi}{2}\right) \times(0,2 \pi)
$$

## Discrete surfaces

$$
\begin{gathered}
f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}, \quad\left(i_{1}, \ldots, i_{d}\right) \mapsto f\left(i_{1}, \ldots, i_{d}\right)=f_{i_{1}, \ldots, i_{d}} \\
\left(f_{i_{1}, \ldots, i_{j}+1, \ldots, i_{k} \ldots i_{d}}-f_{i_{1}, \ldots, i_{j}, \ldots, i_{k} \ldots, i_{d}}\right) \times\left(f_{i_{1}, \ldots, j_{j}, \ldots, i_{k}+\ldots, \ldots, i_{d}}-f_{i_{1}, \ldots, \ldots, \ldots, i_{k} \ldots, i_{d}}\right) \neq 0 \\
f: \mathbb{Z}^{2} \rightarrow \mathbb{R}^{3}, \quad(i, j) \mapsto f(i, j)=f_{i, j} \\
\left(f_{i+1, j}-f_{i j}\right) \times\left(f_{i, j+1}-f_{i j}\right) \neq 0
\end{gathered}
$$

## Shift notation

- $\tau_{j}$ : shift in $j$-th coordinate direction, that is, $\tau_{j} f_{i_{1}, \ldots, i_{j}, \ldots, i_{d}}=f_{i_{1}, \ldots, i_{j}+1, \ldots, i_{d}}$
- write $f, f_{1}, f_{2}, f_{12}$ etc. instead of $f_{i j}, \tau_{1} f_{i j}, \tau_{2} f_{i j}, \tau_{1} \tau_{2} f_{i j}$ etc., for example $\left(f_{i}-f\right) \times\left(f_{j}-f\right) \neq 0$


## Surface curves

$$
\begin{gathered}
\gamma(t)=f(u(t), v(t)) \\
\dot{\gamma}(t)=\frac{\mathrm{d} \gamma}{\mathrm{~d} t}=\frac{\partial f}{\partial u} \frac{\mathrm{~d} u}{\mathrm{~d} t}+\frac{\partial f}{\partial v} \frac{\mathrm{~d} v}{\mathrm{~d} t}
\end{gathered}
$$



- tangents of all surface curve through a fixed surface point $f$ lie in the plane through $f$ and parallel to $\frac{\partial f}{\partial u}$ and $\frac{\partial f}{\partial v}$
- tangent plane $T$ is parallel to $\frac{\partial f}{\partial u}$ and $\frac{\partial f}{\partial v}$
- surface normal $N$ is parallel to $n=\frac{\partial f}{\partial u} \times \frac{\partial f}{\partial v}$


## Conjugate parametrization

## Definition

A surface parametrization $f(u, v)$ is called a conjugate parametrization if
$f_{u}=\frac{\partial f}{\partial u}, f_{v}=\frac{\partial f}{\partial v}$, and $f_{u v}=\frac{\partial^{2} f}{\partial u \partial v}$

are linearly dependent for
every pair $(u, v)$.

- *invariant under projective transformations
- *tangents of parameter lines of one kind along one parameter line of the other kind form a torse
- conjugate directions belong to light ray and corresponding shadow boundary
- conjugate directions with respect to Dupin indicatrix


## Examples

## Example

Show that the surface parametrization

$$
f(u, v)=\frac{1}{\cos u+\cos v-2}\left(\begin{array}{c}
\sin u-\sin v \\
\sin u+\sin v \\
\cos v-\cos u
\end{array}\right)
$$

is a conjugate parametrization.

## Solution

```
with(LinearAlgebra):
F := 1/(cos(u)+\operatorname{cos}(v)-2) *
    Vector([sin(u)-sin(v), sin(u)+sin(v), cos(v)-cos(u)]):
Fu := map(diff, F, u): Fv := map(diff, F, v):
Fuv := map(diff, Fu, v):
Rank(Matrix([Fu, Fv, Fuv]));
```


## Examples

## Example

Assume that the rational bi-quadratic tensor-product Bézier-surface

$$
f(u, v)=f(u, v)=\frac{\sum_{i=0}^{2} \sum_{j=0}^{2} w_{i j} p_{i j} B_{i}^{2}(u) B_{j}^{2}(v)}{\sum_{i=0}^{2} \sum_{j=0}^{2} w_{i j} B_{i}^{2}(u) B_{j}^{2}(v)}
$$

defines a conjugate parametrization. Show that in this case the four sets of control points

$$
\begin{array}{ll}
\left\{p_{00}, p_{01}, p_{11}, p_{10}\right\}, & \left\{p_{01}, p_{02}, p_{12}, p_{11}\right\} \\
\left\{p_{10}, p_{11}, p_{21}, p_{20}\right\}, & \left\{p_{11}, p_{12}, p_{22}, p_{21}\right\}
\end{array}
$$

are necessarily co-planar.

## Examples



Solution

- $w_{00} f_{u}(0,0)=2 w_{10}\left(p_{10}-p_{00}\right)$, $w_{00} f_{v}(0,0)=2 w_{01}\left(p_{01}-p_{00}\right)$
- $4 w_{00}^{2} f_{u v}(0,0)=$
$w_{00} w_{11}\left(p_{11}-p_{00}\right)-w_{01} w_{10}\left(\left(p_{01}-p_{00}\right)+\left(p_{10}-p_{00}\right)\right)$


## Discrete conjugate nets

## Definition

A discrete surface $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}$ is called a discrete conjugate surface (or a conjugate net), if every elementary quadrilateral is planar, that is, if the three vectors

$$
f_{i}-f, \quad f_{j}-f, \quad f_{i j}-f
$$

are linearly dependent for
 $1 \leqslant i<j \leqslant d$.

- *invariant under projective transformations
- *edges in one net direction along thread in other net direction form a discrete torse


## Analytic description of conjugate nets

$$
f_{i j}=f+c_{j i}\left(f_{i}-f\right)+c_{i j}\left(f_{j}-f\right), \quad c_{j i}, c_{i j} \in \mathbb{R}
$$

Construction of a conjugate net $f$ from

1. values of $f$ on the coordinate axes of $\mathbb{Z}^{d}$ and
2. $d(d-1)$ scalar functions $c_{j i}, c_{i j}: \mathbb{Z}^{d} \rightarrow \mathbb{R}$

## Example

For which values of $c_{j i}$ and $c_{i j}$ is the quadrilateral $f f_{1} f_{12} f_{2}$

1. convex,
2. embedded?

## Solution

By an affine transformation, the situation is equivalent to

$$
f=(0,0), \quad f_{i}=(1,0), \quad f_{j}=(0,1) .
$$

Then the fourth vertex is $f_{i j}=\left(c_{j i}, c_{i j}\right)$. The quadrilateral is

- convex if $c_{i j}, c_{i j} \geqslant 0$ and

$$
c_{j i}+c_{i j} \geqslant 1
$$

- embedded if
- $c_{j i}+c_{i j}>1$ or
- $c_{j i}, c_{i j}>0$ or
- $c_{j i}=0, c_{i j} \geqslant 1$ or
- $c_{i j}=0, c_{j i} \geqslant 1$ or
- $c_{j i}, c_{i j}<0$.

convex embedded


## The basic 3D system

## Theorem

Given seven vertices $f, f_{1}, f_{2}, f_{3}, f_{12}, f_{13}$, and $f_{23}$ such that each quadruple $f f_{i} f_{j} f_{i j}$ is planar there exists a unique point $f_{i j k}$ such that each quadruple $f_{i} f_{i j} f_{i k} f_{i j k}$ is planar.

## Proof.

- The initially given vertices lie in a three-space.
- The point $f_{123}$ is obtained as intersection of three planes in this three-space.


## 3D consistency of a 2D system



## 4D consistency of a 3D system



## 4D consistency of a 3D system



## 4D consistency of a 3D system



## 4D consistency of a 3D system



## 4D consistency of a 3D system



## 4D consistency of conjugate nets

Theorem
The 3D system governing discrete conjugate nets is $4 D$ consistent.
Proof.
More-dimensional geometry.
Corollary
The 3D system governing discrete conjugate nets is $n D$ consistent.
Proof.
General result of combinatorial nature on 4D consistent 3D systems.

## Quadric restriction of conjugate nets

## Theorem

Given seven vertices $f, f_{1}, f_{2}, f_{3}, f_{12}, f_{13}$, and $f_{23}$ on a quadric $Q$ such that each quadruple $f f_{i} f_{j} f_{i j}$ is planar, there exists a unique point $f_{i j k} \in Q$ such that each quadruple $f_{i} f_{i j} f_{i k} f_{i j k}$ is planar.

## circular-net

## Lemma

Given seven generic points $f, f_{1}, f_{2}, f_{3}, f_{12}, f_{13}, f_{23}$ in three space there exists an eighth point $f_{123}$ such that any quadric through $f, f_{1}, f_{2}, f_{3}$, $f_{12}, f_{13}, f_{23}$ also contains $f_{123}$.
Proof.

- Quadric equation: $[1, x] \cdot Q \cdot[1, x]=0$ with $Q \in \mathbb{R}^{4 \times 4}$, symmetric, unique up to constant factor
- Quadrics through $f, \ldots f_{23}: \lambda_{1} Q_{1}+\lambda_{2} Q_{2}+\lambda_{3} Q_{3}=0$ (solution system of seven linear homogeneous equations)
- $f_{123}=Q_{1} \cap Q_{2} \cap Q_{3} \backslash\left\{f_{1} \ldots f_{23}\right\}$


## Quadric restriction of conjugate nets

Theorem
Given seven vertices $f, f_{1}, f_{2}, f_{3}, f_{12}, f_{13}$, and $f_{23}$ on a quadric $Q$ such that each quadruple $f f_{i} f_{j} f_{i j}$ is planar, there exists a unique point $f_{i j k} \in Q$ such that each quadruple $f_{i} f_{i j} f_{i k} f_{i j k}$ is planar.

Proof.

- The 3D systems determines $f_{i j k}$ uniquely.
- The pair of planes $f \vee f_{i} \vee f_{j} \vee f_{i j}$ and $f_{k} \vee f_{i k} \vee f_{j k}$ is a (degenerate) quadric through the initially given points.
- Three quadrics of this type intersect in $f_{i j k}$.


## The meaning of quadric restriction

Conjugate nets in quadric models of geometries:

- line geometry (Plücker quadric)
- geometry of SE(3) (Study quadric)
- geometry of oriented spheres (Lie quadric)

Conjugate nets in intersection of quadrics:

- geometry of SE(3) (intersection of six quadrics in $\mathbb{R}^{12}$ )

Specializations of conjugate nets:

- circular nets
- . .


## The meaning of 3D consistency



## Literature

Q R. Sauer
Differenzengeometrie
Springer (1970)
A. I. Bobenko, Yu. B. Suris

Discrete Differential Geometrie. Integrable Structure American Mathematical Society (2008)

## Numeric computation of conjugate nets

Contradicting aims

- planarity
- fairness
- closeness to given surface

Planarity criteria

- $\alpha+\beta+\gamma+\delta-2 \pi=0$
(planar and convex)
- distance of diagonals
- $\operatorname{det}\left(a, a_{j}, b\right)=\cdots=0$, (planar, avoid singularities)
- minimize a linear combination of
- fairness functional and
- closeness functional subject to planarity constraints



## Literature

围 Liu Y., Pottmann H., Wallner J., Yang Y.-L., Wang W. Geometric Modeling with Conical and Developable Surfaces ACM Transactions on Graphics, vol. 25, no. 3, 681-689.

围 Zadravec M., Schiftner A., Wallner J.
Designing quad-dominant meshes with planar faces. Computer Graphics Forum 29/5 (2010), Proc. Symp. Geometry Processing, to appear.

## Asymptotic parametrization

## Definition

A surface parametrization $f(u, v)$ is called an asymptotic parametrization if

$$
\frac{\partial f}{\partial u}, \frac{\partial f}{\partial v}, \frac{\partial^{2} f}{\partial u^{2}} \quad \text { and } \quad \frac{\partial f}{\partial u}, \frac{\partial f}{\partial v}, \frac{\partial^{2} f}{\partial v^{2}}
$$

are linearly dependent for every pair $(u, v)$.
Asymptotic lines

- exist only on surfaces with hyperbolic curvature
- *osculating plane of parameter lines is tangent to surface (rectifying plane contains surface normal)
- intersection curve of surface and rectifying plane of parameter lines has an inflection point
- invariant under projective transformations


## An Example

## Example

Show that the surface parametrization

$$
f(u, v)=\left(\begin{array}{c}
u \\
v \\
u v
\end{array}\right)
$$

is an asymptotic parametrization.

## Solution

We compute the partial derivative vectors:

$$
f_{u}=(1,0, v), \quad f_{v}=(0,1, u), \quad f_{u u}=f_{v v}=(0,0,0)
$$

Obviously, $f_{u u}$ and $f_{v v}$ are linearly dependent from $f_{u}$ and $f_{v}$.

## A pseudosphere



圆 Wunderlich W.
Zur Differenzengeometrie der Flächen konstanter negativer Krümmung Österreich. Akad. Wiss. Math.-Naturwiss. Kl. S.-B.
II, vol. 160, no. 2, 39-77, 1951.

## Discrete asymptotic nets

## Definition

A discrete surface $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{3}$ is called a discrete asymptotic surface (or an asymptotic net), if there exists a plane through $f$ that contains all vectors

$$
f_{i}-f, \quad f_{-i}-f
$$

for $1 \leqslant i \leqslant d$ (planar "vertex stars").

- well-defined tangent plane $T$ and surface normal $N$ at every vertex $f$
- discrete partial derivative vector $\left(f_{i}-f\right)+\left(f-f_{-i}\right)$ is parallel to $T$


## Examples

A sportive example
http://www.flickr.com/photos/laffy4k/202536862/
http://www.flickr.com/photos/bekahstargazing/436888403/
http://www.flickr.com/photos/nataliefranke/2785575144/
A floristic example
blumenampel-1.jpg blumenampel-2.jpg
An architectural example
http://www.flickr.com/photos/preef/4610086160/

## Properties of asymptotic nets

- *invariant under projective transformations
- *asymptotic lines have osculating planes tangent to the surface

Asymptotic nets in higher dimension

- straightforward extension to maps $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}$
- nonetheless only asymptotic nets in a three-space


## Construction of 2D asymptotic nets

- Prescribe values of $f$ on coordinate axes such that all vectors

$$
\tau_{i} f_{0,0}-f_{0,0}, \quad i \in\{1,2\}
$$

are parallel to a plane.

- $f_{1,1}$ lies in the intersection of the two planes

$$
f_{0,0} \vee f_{1,0} \vee f_{2,0} \quad \text { and } \quad f_{0,0} \vee f_{0,1} \vee f_{0,2}
$$

(one degree of freedom)

- inductively construct remaining values of $f$ (one degree of freedom per vertex)


## Construction of asymptotic nets in dimension three

- Prescribe values of $f$ on coordinate axes such that all vectors

$$
\tau_{i} f_{0,0,0}-f_{0,0,0}, \quad i \in\{1,2,3\}
$$

are parallel to a plane.

- Complete the points

$$
\tau_{i} \tau_{j} f_{0,0,0}, \quad i, j \in\{1,2,3\} ; i \neq j
$$

(one degree of freedom per vertex).

- three ways to construct $f_{1,1,1}$ from the already constructed values $\Longrightarrow$ three straight lines

Do these lines intersect? Are asymptotic nets governed by a 3D system?

## Möbius tetrahedra

## Definition

Two tetrahedra $a_{0} a_{1} a_{2} a_{3}$ and $b_{0} b_{1} b_{2} b_{3}$ are called Möbius tetrahedra, if

$$
a_{i} \in b_{j} \vee b_{k} \vee b_{l} \quad \text { and } \quad b_{i} \in a_{j} \vee a_{k} \vee a_{l}
$$

for all pairwise different $i, j, k, l \in\{0,1,2,3\}$.
(Points of one tetrahedron lie in corresponding planes of the other tetrahedron.)

## Theorem (Möbius)

Seven of the eight incidence relations ( $*$ ) imply the eighth.

## Möbius tetrahedra

## Proof.

1. Notation: $A_{i}=a_{j} \vee a_{k} \vee a_{l}$, $B_{i}=b_{j} \vee b_{k} \vee b_{l}$
2. Choose $a_{0}, B_{0}$ with $a_{0} \in B_{0}$.
3. Choose $a_{1}, a_{2}, a_{3}$ (general position $) \rightsquigarrow A_{0}, A_{1}, A_{2}, A_{3}$.
4. Choose $b_{1} \in B_{0} \cap A_{1}$, $b_{2} \in B_{0} \cap A_{2}, b_{3} \in B_{0} \cap A_{3} \rightsquigarrow$ $B_{1}=b_{2} \vee b_{3} \vee a_{1}$,
$B_{2}=b_{1} \vee b_{3} \vee a_{2}$,
$B_{3}=b_{1} \vee b_{2} \vee a_{3}$.

5. $b_{0}:=B_{1} \cap B_{2} \cap B_{3}$, Claim: $b_{0} \in A_{0}$ ( $\checkmark$ by Pappus' Theorem).

## Construction of asymptotic nets in dimension three (II)

- Asymptotic net $\sim$ pairs $(f, T)$ of points $f$ and planes $T$ with $f \in T$; defining property

$$
f \in \tau_{i} T \quad \text { and } \quad \tau_{i} f \in T
$$

- Partition the vertices of the elementary hexahedron of an asymptotic net into two vertex sets of tetrahedra:

$$
\begin{aligned}
& a_{0}=f_{0,0,0}, a_{1}=f_{1,1,0}, a_{2}=f_{1,0,1}, a_{3}=f_{0,1,1} \\
& b_{0}=f_{1,1,1}, b_{1}=f_{0,0,1}, b_{2}=f_{0,1,0}, b_{3}=f_{1,0,0}
\end{aligned}
$$

- Construction of the vertices $f_{i j k}$ with $(i, j, k) \neq(1,1,1)$ yields the configuration of Möbius' Theorem
$\Longrightarrow$ construction of $f_{111}$ without contradiction.


## Analytic description of asymptotic nets

Asymptotic net: $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{3}$
Lelieuvre vector field: $n: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{3}$ such that

$$
\begin{aligned}
& \text { 1. } n \perp T \text { and } \\
& \text { 2. } f_{i}-f=n_{i} \times n
\end{aligned}
$$

- vector $n_{i}$ can be constructed uniquely from $f, n, f_{i}$ (three linear equations)
- vector $n_{i j}$ can be constructed via
- $f, n, f_{i} \rightsquigarrow n_{i} ; f_{i j} \rightsquigarrow n_{i j}$
- $f, n, f_{j} \rightsquigarrow n_{j} ; f_{i j} \rightsquigarrow n_{i j}$

Do these values coincide?

## An auxiliary result

## Lemma (Product formula)

Consider a skew quadrilateral $f, f_{i}, f_{i j}$, $f_{j}$ and vectors $n, n_{i}, n_{i j}, n_{j}$ such that

$$
f_{i}-f=\alpha n_{i} \times n, \quad f_{j}-f=\beta n_{j} \times n,
$$

$$
f_{i j}-f_{j}=\alpha_{j} n_{j} \times n_{j}, \quad f_{i j}-f_{i}=\beta_{i} n_{i j} \times n_{i}
$$

Then $\alpha \alpha_{j}=\beta \beta_{i}$.


## Proof.

- $\left(f_{i}-f\right)^{\mathrm{T}} \cdot n_{j}=\alpha\left(n_{i} \times n\right)^{\mathrm{T}} \cdot n_{j}=-\alpha\left(n_{j} \times n\right)^{\mathrm{T}} \cdot n_{i}$
- $\left(f_{j}-f\right)^{\mathrm{T}} \cdot n_{i}=\beta\left(n_{j} \times n\right)^{\mathrm{T}} \cdot n_{i}$
$--\frac{\alpha}{\beta}=\frac{\left(f_{i}-f\right)^{\mathrm{T}} \cdot n_{j}}{\left(f_{j}-f\right)^{\mathrm{T}} \cdot n_{i}}=\frac{\left(f_{i}-f+f-f_{j}\right)^{\mathrm{T}} \cdot n_{j}}{\left(f_{j}-f+f-f_{i}\right)^{\mathrm{T}} \cdot n_{i}}=\frac{\left(f_{i}-f_{j}\right)^{\mathrm{T}} \cdot n_{j}}{\left(f_{j}-f_{i}\right)^{\mathrm{T}} \cdot n_{i}}$


## An auxiliary result

## Lemma (Product formula)

Consider a skew quadrilateral $f, f_{i}, f_{i j}$, $f_{j}$ and vectors $n, n_{i}, n_{i j}, n_{j}$ such that

$$
f_{i}-f=\alpha n_{i} \times n, \quad f_{j}-f=\beta n_{j} \times n,
$$

$$
f_{i j}-f_{j}=\alpha_{j} n_{j} \times n_{j}, \quad f_{i j}-f_{i}=\beta_{i} n_{i j} \times n_{i}
$$

Then $\alpha \alpha_{j}=\beta \beta_{i}$.


Proof.

$$
\begin{aligned}
& -\frac{\alpha}{\beta}=\frac{\left(f_{i}-f\right)^{\mathrm{T}} \cdot n_{j}}{\left(f_{j}-f\right)^{\mathrm{T}} \cdot n_{i}}=\frac{\left(f_{i}-f+f-f_{j}\right)^{\mathrm{T}} \cdot n_{j}}{\left(f_{j}-f+f-f_{i}\right)^{\mathrm{T}} \cdot n_{i}}=\frac{\left(f_{i}-f_{j}\right)^{\mathrm{T}} \cdot n_{j}}{\left(f_{j}-f_{i}\right)^{\mathrm{T}} \cdot n_{i}} \\
& >-\frac{\alpha_{j}}{\beta_{i}}=\cdots=\frac{\left(f_{i}-f_{j}\right)^{\mathrm{T}} \cdot n_{i}}{\left(f_{i}-f_{j}\right)^{\mathrm{T}} \cdot n_{j}} \\
& \Rightarrow \Longrightarrow \frac{\alpha}{\beta}=\frac{\beta_{i}}{\alpha_{j}}
\end{aligned}
$$

## Existence and uniqueness

## Theorem

The Lelieuvre normal vector field $n$ of an asymptotic net $f$ is uniquely determined by its value at one point.

Proof.
Uniqueness $\checkmark$

## Existence

- Product formula for normal vector fields: $\alpha \alpha_{j}=\beta \beta_{i}$.
- Three of the values $\alpha, \alpha_{j}, \beta, \beta_{i}$ equal $1 \Longrightarrow$ all four values equal 1.
- The Lelieuvre normal vector field is characterized by $\alpha=\alpha_{j}=\beta=\beta_{i}=1$.
- Both construction of $n_{i j}$ result in the same value.


## Relation between two Lelieuvre normal vector fields

Theorem
Suppose that $n$ and $n^{\prime}$ are two Lelieuvre normal vector fields to the same asymptotic net. Then there exists a value $\alpha \in \mathbb{R}$ such that

$$
n(z)= \begin{cases}\alpha n(z) & \text { if } z_{1}+\cdots+z_{d} \text { is even } \\ \alpha^{-1} n(z) & \text { if } z_{1}+\cdots+z_{d} \text { is odd } .\end{cases}
$$

Proof. $\checkmark$

## The discrete surface of Lelieuvre normals

What are the properties of the discrete net $n: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{3}$ ?

- $f_{i j}-f=f_{i j}-f_{i}+f_{i}-f=n_{i j} \times n_{i}+n_{i} \times n$
- $f_{i j}-f=f_{i j}-f_{j}+f_{j}-f=n_{i j} \times n_{j}+n_{j} \times n$
$\Rightarrow \Longrightarrow\left(n_{i j}-n\right) \times\left(n_{i}-n_{j}\right)=0$
- $\Longrightarrow n_{i j}-n=a_{i j}\left(n_{j}-n_{i}\right)$ where $a_{i j} \in \mathbb{R}$

Conclusion:

- The net $n: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{3}$ is conjugate.
- Every fundamental quadrilateral has parallel diagonals (this is called a "T-net").


## T-nets

## Defining equation:

$$
y_{i j}-y=a_{i j}\left(y_{j}-y_{i}\right) \quad \text { where } \quad a_{i j} \in \mathbb{R}
$$

- $a_{i j}=-a_{j i}$
- $y_{i j}-y=\left(1+c_{j i}\right)\left(y_{i}-y\right)+\left(1+c_{i j}\right)\left(y_{j}-y\right) \Longrightarrow$
- $c_{i j}+c_{j i}+2=0$ (T-net condition)
- $a_{i j}=c_{i j}+1$ (relation between coefficients)


## Elementary hexahedra of T-nets

## Theorem

Consider seven points $y, y_{1}, y_{2}, y_{3}, y_{12}, y_{13}, y_{23}$ of a combinatorial cube such that the diagonals of

$$
y y_{1} y_{12} y_{2}, \quad y y_{1} y_{13} y_{3}, \quad \text { and } y y_{2} y_{23} y_{3}
$$

are parallel. Then there exists a unique point $y_{123}$ such that also the diagonals of

$$
y_{1} y_{12} y_{123} y_{13}, \quad y_{2} y_{12} y_{123} y_{23}, \quad \text { and } y_{3} y_{13} y_{123} y_{23}
$$

are parallel.
Corollary
T-nets are described by a 3D system. They are $n D$ consistent.

## Elementary hexahedra of T-nets

## Proof.

- $y_{i j}-y=a_{i j}\left(y_{j}-y_{i}\right) \Longrightarrow$

$$
\tau_{i} y_{j k}=\left(1+\left(\tau_{i} a_{j k}\right)\left(a_{i j}+a_{k i}\right)\right) y_{i}-\left(\tau_{i} a_{j k}\right) a_{i j} y_{j}-\left(\tau_{i} a_{j k}\right) a_{k i} y_{k}
$$

- Six linear conditions for three unknowns $\tau_{i} a_{j k}$ :

$$
\begin{aligned}
& 1+\left(\tau_{1} a_{23}\right)\left(a_{12}+a_{31}\right)=-\left(\tau_{2} a_{31}\right) a_{12}=-\left(\tau_{3} a_{12}\right) a_{31} \\
& 1+\left(\tau_{2} a_{31}\right)\left(a_{23}+a_{12}\right)=-\left(\tau_{3} a_{12}\right) a_{23}=-\left(\tau_{1} a_{23}\right) a_{12} \\
& 1+\left(\tau_{2} a_{30}\right)\left(a_{23}+a_{02}\right)=-\left(\tau_{3} a_{02}\right) a_{23}=-\left(\tau_{0} a_{23}\right) a_{02}
\end{aligned}
$$

- Unique solution:

$$
\frac{\tau_{1} a_{23}}{a_{23}}=\frac{\tau_{2} a_{31}}{a_{31}}=\frac{\tau_{3} a_{12}}{a_{12}}=\frac{1}{a_{12} a_{23}+a_{23} a_{31}+a_{31} a_{12}}
$$

## Asymptotic nets from T-nets

Theorem
An asymptotic net is uniquely defined (up to translation) by a Lelieuvre normal vector field (a T-net).

## Corollary

Asymptotic nets are $n D$ consistent.
Question: How to construct an asymptotic net from a given T-net $n$ ?

## Discrete one forms

- graph $G$ with vertex set $V$, set of directed edges $\vec{E}$
- vector space $W$

Definition (discrete additive one-form)

- $p: \vec{E} \rightarrow W$ is a discrete additive one-form if $p(-e)=-p(e)$.
- $p$ is exact if $\sum_{e \in Z} p(e)=0$ for every cycle $Z$ of directed edges.

Example: $p(e)=e$.
Definition (discrete multiplicative one-form)

- $q: \vec{E} \rightarrow \mathbb{R} \backslash 0$ is a discrete multiplicative one-form if $q(-e)=1 / q(e)$.
- $q$ is exact if $\prod_{e \in Z} q(e)=1$ for every cycle $Z$ of directed edges.


## Integration of exact forms

## Theorem

Given the exact additive discrete one form $p: \vec{E} \rightarrow W$ there exists a function $f: V \rightarrow W$ such that $p(e)=f(y)-f(x)$ for any $e=(x, y)$ in $\vec{E}$. The function $f$ is defined up to an additive constant.
Proof. $\checkmark$

## Theorem

Given the exact multiplicative discrete one form $q: \vec{E} \rightarrow \mathbb{R} \backslash 0$ there exists a function $v: V \rightarrow \mathbb{R} \backslash 0$ such that $q(e)=v(y) / v(x)$ for any $e=(x, y)$ in $\vec{E}$. The function $v$ is defined up to an additive constant.

## Integration of exact forms

Theorem
Given the exact additive discrete one form $p: \vec{E} \rightarrow W$ there exists a function $f: V \rightarrow W$ such that $p(e)=f(y)-f(x)$ for any $e=(x, y)$ in $\vec{E}$. The function $f$ is defined up to an additive constant. Proof. $\checkmark$

Question: How to construct an asymptotic net from a given T-net $n$ ?

Answer: Integrate the exact one form $p(i, j)=n_{i} \times n_{j}$.

## Ruled surfaces and torses

$\mathcal{L}^{n} \ldots$ set of lines in $\mathbb{R}^{( }{ }^{n}$ (typically $n=3$ )
Definition
A ruled surface is a (sufficiently regular) map $\ell: \mathbb{R} \rightarrow \mathcal{L}^{n}$.
Definition
A discrete ruled surface is a map $\ell: \mathbb{Z} \rightarrow \mathcal{L}^{n}$ such that
$\ell \cap \ell_{i}=\varnothing$.

## Definition

A torse is a map $\ell: \mathbb{R} \rightarrow \mathcal{L}^{n}$ such that all image lines are tangent to a (sufficiently regular) curve.

## Definition

A discrete torse is a map $\ell: \mathbb{Z} \rightarrow \mathcal{L}^{n}$ such that $\ell \cap \ell_{i} \neq \varnothing$.
$\Longrightarrow$ existence of polygon of regression, osculating planes etc.

## Smooth line congruences

## Definition

A line congruence is a (sufficiently regular) map $\ell: \mathbb{R}^{2} \rightarrow \mathcal{L}^{n}$.

## Examples

- normal congruence of a smooth surface: $f(u, v)+\lambda n(u, v)$ where $n=f_{u} \times f_{v}$.
- set of transversals of two skew lines
- sets of light rays in geometrical optics


## Discrete line congruences

## Definition

A discrete line congruence is a map $\ell: \mathbb{Z}^{d} \rightarrow \mathcal{L}^{n}$ such that any two neighbouring lines $\ell$ and $\ell_{i}$ intersect.

- smooth line congruences admit special parametrizations $\rightsquigarrow$ different discretizations conceivable
- discretize definition considers only parametrization "along torses"


## Construction of discrete line congruences

$d=2: \checkmark$
$d=3$ : The completion of an elementary hexahedron from seven lines $\ell, \ell_{1}, \ell_{2}, \ell_{3}, \ell_{12}, \ell_{13}, \ell_{23}$ is possible and unique (3D system).
$d=4$ : The completion of an elementary hypercube from 15 lines $\ell, \ell_{i}, \ell_{i j}, \ell_{i j k}$ is possible and unique (4D consistent).
$d>4 n \mathrm{D}$ consistent

## Discrete line congruences and conjugate nets

## Definition

The $i$-th focal net of a discrete line congruence $\ell: \mathbb{Z}^{d} \rightarrow \mathcal{L}^{n}$ is defined as $F^{(i)}=\ell \cap \ell_{i}$.

## Theorem

The $i$-th focal net of a discrete line congruence is a discrete conjugate net.

## Theorem

Given a discrete conjugate net $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}$, a discrete line congruence $\ell: \mathbb{Z}^{d} \rightarrow \mathcal{L}^{n}$ with the property $f \in \ell$ is uniquely determined by its values at the coordinate axes in $\mathbb{Z}^{d}$.

## Proof.

Given two lines $\ell_{i}, \ell_{j}$ and a point $f_{i j}$ there exists a unique line $\ell_{i j}$ incident with $f_{i j}$ and concurrent with $\ell_{i}, \ell_{j}$.

## Discrete line congruences and conjugate nets II

## Definition

The $i$-th tangent congruence of a discrete conjugate net $f: \mathbb{Z}^{2} \rightarrow \mathbb{R} \mathbb{P}^{n}$ is defined as $\ell^{(i)}=f \vee f_{i}$.

## Definition

In case of $d=2$ the $i$-th Laplace transform $l^{(i)}$ of a two-dimensional discrete conjugate net is the $j$-th focal congruence of its $i$-th tangent congruence $(i \neq j)$.

Theorem
The Laplace transforms of a discrete conjugate net are discrete conjugate nets.

## Lecture 4: <br> Discrete Curvature Lines

## Curvature line parametrizations



- normal surfaces along parameter lines are torses (infinitesimally neighbouring surface normals along parameter lines intersect)
- $f_{u}, f_{v}$ are tangent to the principal directions
- parameter lines intersect orthogonally


## Discrete curvature line parametrizations

## Neighboring surface normals intersect.

- circular nets
- conical nets
- principal contact element nets
- HR-congruences


## Circular nets

## Definition

A map $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}$ is called a circular net or discrete orthogonal net if all elementary quadrilaterals are circular.


- neighboring circle axes intersect
- discretization of conjugate parametrization


## Algebraic characterization

$$
\begin{gathered}
f_{i j}=f+c_{j i}\left(f_{i}-f\right)+c_{i j}\left(f_{j}-f\right), \quad c_{j i}, c_{i j} \in \mathbb{R} \\
\alpha f+\alpha_{i} f_{i}+\alpha_{j} f_{j}+\alpha_{i j} f_{i j}=0, \quad \alpha+\alpha_{i}+\alpha_{j}+\alpha_{i j}=0 \\
\left(\alpha=1-c_{i j}-c_{j i}, \quad \alpha_{i}=c_{j i}, \quad \alpha_{j}=c_{i j}, \quad \alpha_{i j}=-1\right)
\end{gathered}
$$

Circularity condition:

$$
\alpha\|f\|^{2}+\alpha_{i}\left\|f_{i}\right\|^{2}+\alpha_{j}\left\|f_{j}\right\|^{2}+\alpha_{i j}\left\|f_{i j}\right\|^{2}=0
$$

## Proof.

- $(\star) \Longleftrightarrow \forall m \in \mathbb{R}^{n}:$

$$
\alpha\|f-m\|^{2}+\alpha_{i}\left\|f_{i}-m\right\|^{2}+\alpha_{j}\left\|f_{j}-m\right\|^{2}+\alpha_{i j}\left\|f_{i j}-m\right\|^{2}=0
$$

- Take $m$ as center of circum-circle $C$ of $f, f_{i}, f_{j}$ :

$$
\|f-m\|^{2}=\left\|f_{i}-m\right\|^{2}=\left\|f_{j}-m\right\|^{2}=r^{2}
$$

- $\Longrightarrow\left\|f_{i j}-m\right\|=r^{2} \Longrightarrow f_{i j} \in C$


## Circularity criteria

## Theorem

The four points $a, b, c, d \in \mathbb{R}^{2}$ lie on $a$ circle if and only if opposite angles in the quadrilateral $a b c d$ are supplementary, that is,

$$
\alpha+\gamma=\beta+\delta=\pi
$$

(immediate consequence from the Inscribed-Angle Theorem)


## Circularity criteria

## Theorem

The four points $a, b, c, d \in \mathbb{C}$ lie on a circle (or a straight line) if and only if

$$
\frac{a-b}{b-c} \cdot \frac{c-d}{d-a} \in \mathbb{R}
$$

## Proof.

- Angle between complex numbers equals argument of their ratio: $\varangle(a, b)=\arg (a / b)$
- Two complex numbers $a, b$ have the same or supplimentary argument $\Longleftrightarrow a / b \in \mathbb{R}$.
- ( $\star$ ) equals

$$
\frac{a-b}{c-b}: \frac{a-d}{c-d}
$$

and thus states equality or supplimentary of $\beta$ and $\delta$.

## Circularity criteria

Theorem
The four points $a, b, c, d \in \mathbb{C}$ lie on a circle (or a straight line) if and only if

$$
\frac{a-b}{b-c} \cdot \frac{c-d}{d-a} \in \mathbb{R}
$$

Cross-ratio criterion for circularity:

$$
C R(a, b, c, d)=\frac{a-c}{b-c} \cdot \frac{b-d}{a-d} \in \mathbb{R}
$$

- better known
- more difficult to memorize
- similar proof (use Incident Angle Theorem)


## Circularity criteria

In the following theorem, $a, b, c$, and $d$ are considered as vector valued quaternions; multiplication (not commutative) and inversion are performed in the quaternion division ring.

## Theorem

The four points $a, b, c, d \in \mathbb{R}^{3}$ lie on a circle (or a straight line) if and only if their cross-ratio

$$
\mathrm{CR}(a, b, c, d)=(a-b) \star(b-c)^{-1} \star(c-d) \star(d-a)^{-1}
$$

is real.
Proof. cross-ratio-criterion.mw

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## Two-dimensional circular nets

## Defining data

- values of $f$ on coordinate axes of $\mathbb{Z}^{2}$
- a cross-ratio on each elementary quadrilateral

Shape of the circles
The quadrilateral $a b c d$ is circular and embedded if and only if

$$
\frac{a-b}{b-c} \cdot \frac{c-d}{d-a}<0
$$

Numerical computation
Add circularity condition
$\sum(\alpha+\gamma-\pi)^{2}+\sum(\beta+\delta-\pi)^{2} \rightarrow \min$ to optimization scheme.

## Three-dimensional circular nets

## Theorem

Circular nets are governed by a 3D system.

## Theorem

Given seven vertices $f, f_{1}, f_{2}, f_{3}, f_{12}, f_{13}$, and $f_{23}$ such that each quadruple $f f_{i} f_{j} f_{i j}$ lies on a circle, there exists a unique point $f_{i j k}$ such that each quadruple $f_{i} f_{i j} f_{i k} f_{i j k}$ is a circular quadrilateral.

Proof.

- All initially given vertices lie on a sphere $S$.
- Claim follows from quadric reduction of conjugate nets.

Alternative: Miquel's Six Circles Theorem

## Conical nets

## Definition

A map
$P: \mathbb{Z}^{d} \rightarrow\left\{\right.$ oriented planes in $\left.\mathbb{R}^{3}\right\}$
is called a conical net the four planes $P, P_{i}, P_{i j}, P_{j}$ are tangent to an oriented cone of revolution.


- neighboring cone axes intersect
- discretization of conjugate parametrization


## The Gauss map of conical nets

- Every plane $P$ is described by unit normal $n$ and distance $d$ to the origin.
- The map $n: \mathbb{Z}^{d} \rightarrow S^{2} \subset \mathbb{R}^{3}$ is the Gauss map of the conical net.

Theorem
The Gauss map is circular.

- A conical net is uniquely determined by its Gauss map and the map $d: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{+}$.
- Conicality criterion:

$$
\left(n-n_{i}\right) \star\left(n_{i}-n_{i j}\right)^{-1} \star\left(n_{i j}-n_{j}\right) \star\left(n_{j}-n\right)^{-1} \in \mathbb{R}
$$

## Circular quadrilaterals

Theorem
The composition of the reflections in two intersecting lines is a rotation about the intersection point through twice the angle between the two lines.


## Theorem

The composition of reflections in successive bisector planes of a circular quadrilateral yields the identity.


## Conical nets from circular nets

## Theorem

Given a circular net $f$ there exists a two-parameter variety of conical nets whose face planes are incident with the vertices of $f$. Any such net is uniquely determined by one of its face planes.

## Proof.



- Generate the conical net by successive reflection in the bisector planes of neighboring vertices of $f$.
- This construction produces planes of a conical net and is free of contradictions.


## Circular nets from conical nets

Theorem
Given a conical net P there exists a two-parameter variety of circular nets whose vertices are incident with the face planes of $P$. Any such net is uniquely determined by one of its vertices.


## Proof.

Also the composition of the reflections in successive bisector planes of the face planes of a conical net yields the identity.

## Multidimensional consistency

Theorem
Conical nets are governed by a 3D system. They are nD consistent.
Proof.
The claim follow from the analogous statements about circular nets and the fact that both classes of nets can be generated by the same sequence of reflections.

## Literature

A. I. Bobenko, Yu. B. Suris

Discrete Differential Geometrie. Integrable Structure American Mathematical Society (2008)
E H. Pottmann., J. Wallner
The focal geometry of circular and conical meshes Adv. Comput. Math., vol. 29, no. 3, 249-268, 2008.

## Numerical computation

## Theorem (Lexell; Wallner, Liu, Wang)

Consider four unit vectors $e_{0}, e_{1}, e_{2}, e_{3}$ and denote the angle between $e_{i}$ and $e_{i+1}$ by $\psi_{i, i+1}$. The vectors are the directions of the edges emanating from a vertex in a conical net if and only if

$$
\psi_{01}+\psi_{23}=\psi_{12}+\psi_{31} .
$$

- A complete proof considering all possible cases is not difficult but involved.
- The theorem is actually a statement about spherical quadrilaterals with an in-circle.
- For numerical computation, add conicality condition $\sum\left(\psi_{01}+\psi_{23}-\psi_{12}-\psi_{31}\right)^{2} \rightarrow \min$ to optimization scheme.


## Literature

图 Lexell A. J.
Acta Sc. Imp. Petr. (1781) 6, 89-100.
圊 Wang W., Wallner J., Lie Y.
An Angle Criterion for Conical Mesh Vertices
J. Geom. Graphics (2007) 11:2, 199-208.

## HR-congruences



## Definition

A discrete line congruence $\ell: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{3}$ is called an HR-congruence if the skew quadrilateral consisting of the four lines $\ell, \ell_{i}, \ell_{i j}, \ell_{j}$ lies on a hyperboloid of revolution.

## Theorem

If $p$ is a circular net and $T$ a conical net with $p \in T$, then the normals of $T$ form an HR-congruence.
Proof. Construction by reflection.

## Principal contact element nets

## Definition

An (oriented) contact element is a pair $(p, n)$ consisting of a point $p$ and a unit vector $n$.
Alternatively, think of a contact element as

- a pair $(p, N)$ (point plus oriented line),
- a pair $(p, T)$ (point plus oriented tangent plane).


## Definition

A principle contact element net is a map

$$
(p, n): \mathbb{Z}^{d} \rightarrow\{\text { space of oriented contact elements }\}
$$

such that any two neighboring contact elements have a common tangent sphere.

## Properties of principal contact element nets



- The normals of neighboring contact elements intersect in the center of the tangent sphere (curvature line discretization).
- Neighboring contact elements have a unique plane of symmetry.


## Relation to circular and conical nets

## Theorem

Iff is a circular net and $T$ a conical net such that $f \in T$, then $(f, T)$ is a principal contact element net.


## Proof.

Due to the construction by reflections, the intersection points of the plane normals are at the same (oriented distance) from the points of tangency.

## Relation to circular and conical nets

Theorem
If $(p, T)$ is a principal contact element net with face planes $T$, then $p$ is a circular net and $T$ is a conical net.

## Proof.

- Opposite contact elements of an elementary quadrilateral correspond, in two ways, in the composition of two reflections in planes of symmetry.
- Opposite contact elements correspond in two rotations.
- Opposite contact elements have skew normals $\Longrightarrow$ the two rotations are actually identical.
- All four planes of symmetry intersect in a common line and the composition of reflections yields the identity.


## Lecture 5: <br> Parallel Nets, Offset Nets and Curvature

## Parallel nets

## Definition

Let $f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}$ be a conjugate net. A conjugate net $f^{+}: \mathbb{Z}^{\rightarrow} \mathbb{R}^{n}$ is called a parallel net (or a Combescure transform of $f$ ) if corresponding edges are parallel.

Remark
The theory of parallel nets and offset nets as presented below extends to quad meshes of arbitrary combinatorics.


## Parallel nets and line congruences

Given are a conjugate net $f$ and a parallel net $f^{+}$:
$\Longrightarrow \quad \ell=f \vee f^{+}$is a discrete line congruence

Given are a conjugate net $f$ and a discrete line congruence $\ell$ with $f \in \ell$ :
$\Longrightarrow$ There exists a one-parameter family $f^{+}$of parallel nets with $f^{+} \in \ell$.
$\Longrightarrow f^{+}$is uniquely determined by its value at one point.

## Offset nets

## Given:

- conjugate net $f$
- parallel net $f^{+}$


## Definition

A parallel net $f^{+}$is called a vertex/face/edge offset net if corresponding vertices/faces/edges are at constant distance $d$.

## The vector space of parallel nets

## Theorem

All conjugate nets parallel to a given conjugate net form a vector space over $\mathbb{R}$ where addition and multiplication are defined vertex-wise:

$$
\begin{array}{rlrl}
\lambda f: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}, & & i \mapsto \lambda f(i), \\
f+f^{+}: \mathbb{Z}^{d} \rightarrow \mathbb{R}^{n}, & i \mapsto f(i)+f^{+}(i)
\end{array}
$$

## Definition

Let $f$ and $f^{+}$be a pair of offset nets at constant distance $d$. Then the Gauss image of $f^{+}$with respect to $f$ is defined as

$$
s=\frac{1}{d}\left(f^{+}-f\right) .
$$

## The smooth Gauss map for curves



- curvature $\approx$ ratio of arc-lengths of Gauss image and curve


## The smooth Gauss map for surfaces

## Definition

Given a smooth surface $M$, denote by $n_{p}$ the oriented unit normal in $p \in M$. The Gauss map of $M$ is the map

$$
n: M \rightarrow S^{2}, \quad p \mapsto n_{p} .
$$



## The smooth Gauss map for surfaces

## Definition

Given a smooth surface $M$, denote by $n_{p}$ the oriented unit normal in $p \in M$. The Gauss map of $M$ is the map

$$
n: M \rightarrow S^{2}, \quad p \mapsto n_{p}
$$

## Properties:

- closely related to surface curvatures
- negative derivative $-\mathrm{d} n: T_{p}(M) \rightarrow T_{n_{p}}\left(S^{2}\right)$ is called the shape operator


## The Gauss image of offset nets



## Theorem

The Gauss image of a vertex/face/edge offset net is a net

- whose vertices are contained in $S^{d}$,
- whose faces circumscribe $S^{d}$,
- whose edges are tangent to $S^{d}$.


## Characterization of offset-nets

## Corollary

A conjugate net $f$ admits a vertex offset net $f^{+}$if and only if it is circular.

Proof. Assume a vertex offset $f^{+}$exists $\Longrightarrow$ circular Gauss image $\Longrightarrow$ original net is circular (angle criterion for circularity).

Construction of vertex offset nets:
Assume $f$ is circular:

1. Prescribe one vertex of $f^{+}$
2. Construct Gauss image from one vertex and known edge directions (unambiguous; no contradictions by circularity).
3. Construct $f^{+}$from the Gauss image (unambiguous; no contradictions).

## Characterization of offset-nets

Corollary
A conjugate net $f$ admits a face offset net $f^{+}$if and only if it is conical.
Proof. Assume a face offset $f^{+}$exists $\Longrightarrow$ conical Gauss image $\Longrightarrow$ original net is conical (angle criterion for conicality).

Construction of face offset nets:
Assume $f$ is conical:

1. Prescribe one face of $f^{+}$.
2. Construct other faces by offsetting (unambiguous; no contradictions by conicality).

## Characterization of offset-nets

## Definition

A conjugate net is called a Koebe net, if its edges are tangent to the unit sphere.

Corollary
A conjugate net $f$ admits an edge offset net $f^{+}$if and only if it is parallel to a Koebe net s.
Proof. Construction of $f^{+}$from $f$ and $s$ :

$$
f^{+}=f+d \cdot s
$$

## Offset nets in architecture

- fewer edges for quad dominant meshes
- quadrilateral glass panels are cheaper
- less-steel, more glass
- torsion-free nodes
- existence of face or edge offset meshes

R H. Pottmann, Y. Liu, J. Wallner, A. Bobenko, W. Wang Geometry of multi-layer freeform structures for architecture
ACM Trans. Graphics, vol. 26, no. 3, 1-1, 2007

## Discrete line congruences with offset properties

## Definition

Two discrete line congruences $\ell$ and $\ell^{+}$are called parallel, if corresponding lines are parallel.
They are called offset congruences if corresponding lines are at constant distance as well.

Remark
The edges of an edge-offset net constitute a special example of an offset congruence with planar elementary quadrilaterals.

## Remark

Offset congruences occur in architecture of folded paper strips.

## Application: Design of closed folded strips


http://www.archiwaste.org/?p=1109
Institut für Konstruktion und Gestaltung, Universität Innsbruck:
Rupert Maleczek, Eda Schaur
Archiwaste:
Guillaume Bounoure, Chloe Geneveaux

## Offset congruences

## Theorem

All line congruences parallel to a given discrete line congruence l form a vector space. Addition and multiplication are defined via addition and multiplication of corresponding intersection points.

## Definition

The Gauss image of two offset congruences $\ell$ and $\ell^{+}$at distance $d$ is defined as

$$
s=\frac{1}{d}\left(\ell^{+}-\ell\right) .
$$

Theorem
A discrete line congruence $\ell$ admits an offset congruence if and only if it is parallel and at constant distance to a discrete line congruence whose lines are tangent to the unit sphere $S^{2}$.

## Elementary quadrilaterals of the Gauss image



Problem: Given two tangents $A, B$ of $S^{2}$ find lines $X$ which

1. intersect $A$ and $B$ and
2. are tangent to $S^{2}$.

Solution: The locus of possible points of tangency consists of two circles through $a$ and $b$.

## Bi-arcs in the plane and on the sphere


\& H. Pottmann, J. Wallner
Computational Line Geometry
Springer (2001)
围 H. Stachel, W. Fuhs
Circular pipe-connections
Computers \& Graphics 12 (1988), 53-57.

## Elementary quadrilaterals of the Gauss image

## Theorem

Let s be the Gauss image of a pair of offset congruences. An elementary quadrilateral of $s$ is either

1. the elementary quadrilateral of an HR-congruence or
2. something different (yet unnamed)

## Remark

The geometry of offset congruences and metric aspects of discrete line geometry are open research questions.

## Curvature of a smooth curve

$$
\begin{gathered}
\gamma: I \subset \mathbb{R} \rightarrow \mathbb{R}^{3}, \quad t \mapsto \gamma(t), \\
\varkappa(t)=\frac{\|\dot{\gamma}(t) \times \ddot{\gamma}(t)\|}{\|\dot{\gamma}(t)\|^{3}}, \\
l(\gamma)=\int_{I}\|\dot{\gamma}(t)\| \mathrm{d} t .
\end{gathered}
$$



- change of tangent direction per arc-length
- inverse radius of optimally approximating circle


## Steiner's formula

- convex curve $\gamma \subset \mathbb{R}^{2}$, arc-length $s$, curvature $\varkappa(s)$
- offset curve $\gamma_{t}$ at distance $t$

$$
l\left(\gamma_{t}\right)=l(\gamma)+t \int_{\gamma} \varkappa(t) \mathrm{d} t
$$



Example: A circle

$$
l\left(\gamma_{t}\right)=2(r+t) \pi=2 r \pi+2 t \pi=l(\gamma)+t \int_{0}^{2 r \pi} r^{-1} \mathrm{~d} \varphi
$$

## Steiner type curvatures in vertices



- Assign curvature to vertices so that Steiner's Theorem remains true.
- The three possibilities are identical up to second order terms:

$$
2 \sin \frac{\varphi}{2}=\varphi+O\left(\varphi^{3}\right), \varphi=\varphi+O\left(\varphi^{3}\right), 2 \tan \frac{\varphi}{2}=\varphi+O\left(\varphi^{3}\right) .
$$

## Curvatures of a smooth surface



## Gaussian curvature as local area distortion



## Gaussian curvature as local area distortion


area $A$


- principal contact element net $(p, n)$
- Gauss image $n$
- discrete Gauss curvature of a face:

$$
K=\frac{A_{0}}{A}
$$

## Local Steiner formula

Smooth surface $f$, offset surface $f_{t}$ at distance $t$ :

$$
\mathrm{d} A\left(f_{t}\right)=\left(1-2 H t+K t^{2}\right) \mathrm{d} A(f)
$$

- ratio of area elements is a quadratic polynomial in the offset distance
- coefficients depend on Gaussian curvature $K$ and mean curvature $H$


## Discretization:

- compare face areas of offset nets
- use coefficients of (hopefully) quadratic polynomials


## Oriented and mixed area

- $n$-gon $\mathcal{P}=\left\langle p_{0}, \ldots, p_{n-1}\right\rangle \subset \mathbb{R}^{2}$
- oriented area

$$
\begin{aligned}
A(\mathcal{P}) & \left.=\frac{1}{2} \sum_{i=0}^{n} \operatorname{det}\left(p_{i}, p_{i+1}\right) \quad \text { (indices modulo } n\right) \\
& =\left(p_{0}, \ldots, p_{n}\right) \cdot \mathbf{A} \cdot\left(p_{0}, \ldots, p_{n}\right)^{\mathrm{T}} \quad\left(\text { quadratic form in } \mathbb{R}^{2 n}\right)
\end{aligned}
$$

- associated symmetric bilinear form

$$
A(\mathcal{P}, Q)=\left(p_{0}, \ldots, p_{n}\right) \cdot \mathbf{A} \cdot\left(q_{0}, \ldots, q_{n}\right)^{\mathrm{T}}
$$

## Remark

If $P$ and $Q$ are parallel, positively oriented convex polygons then $A(P, Q)$ equals the mixed area (known from convex geometry) of $P$ and $Q$.

## Discrete Steiner formula

- principal contact element net $(f, n)$
- offset net $f_{t}=f+t n$
- corresponding faces $F, F_{t}, N$

$$
\begin{aligned}
A\left(F_{t}\right)= & A(F+t N)= \\
& A(F)+2 t A(F, N)+t^{2} A(N)=\left(1-2 t H+t^{2} K\right) A(F)
\end{aligned}
$$

where

$$
H=-\frac{A(F, S)}{A(F)}, \quad K=\frac{A(S)}{A(F)}
$$

(discrete Gaussian and mean curvature associated to faces)

## Pseudospherical principal contact element nets

Theorem
$\left(f_{0}, n_{0}\right),\left(f_{1}, n_{1}\right),\left(f_{2}, n_{2}\right)$ of an elementary quadrilateral in a principal contact element net, show that there exists precisely one vertex $\left(f_{3}, n_{3}\right)$ such that the Gaussian curvature attains a given value $K$.

- $f_{3}$ is constrained to circle, $n_{3}$ is found by reflection $\rightsquigarrow$ quadratic parametrizations $f_{3}(t)$ and $n_{3}(t)$
- The condition $K \cdot A(F)=A(S)$ is a quadratic polynomial $Q(t)$.
- One of the two zeros of $Q$ is attained for $f_{3}=f_{0}, n_{3}=n_{0}$, the other zero is the sought solution.


## Pseudospherical principal contact element nets

Theorem
$\left(f_{0}, n_{0}\right),\left(f_{1}, n_{1}\right),\left(f_{2}, n_{2}\right)$ of an elementary quadrilateral in a principal contact element net, show that there exists precisely one vertex $\left(f_{3}, n_{3}\right)$ such that the Gaussian curvature attains a given value $K$.

## Corollary

A pseudospherical principal contact element net $(f, n)$ is governed by a 2 D system.

- Kinematic approach, $n \mathrm{D}$ consistency etc. $\rightsquigarrow$ ICGG 2010, CCGG 2010


## Pseudospherical principal contact element nets



## Literature

樯
A. I. Bobenko, H. Pottmann, J. Wallner

A curvature theory for discrete surfaces based on mesh parallelity
Math. Ann., 348:1, 1-24 (2010).
© J.-M. Morvan
Generalized Curvatures
Springer 2008
Q M. Desbrun, E. Grinspun, P. Schröder, M. Wardetzky Discrete Differential Geometry: An Applied Introduction SIGGRAPH Asia 2008 Course Notes

## Lecture 6: <br> Cyclidic Net Parametrization

## Net parametrization

Problem:
Given a discrete structure, find a smooth parametrization that preserves essential properties.

## Examples:

- conjugate parametrization of conjugate nets
- principal parametrization of circular nets
- principal parametrization of planes of conical nets
- principal parametrization of lines of HR-congruence
- ...


## Dupin cyclides



- inversion of torus, revolute cone or revolute cylinder
- curvature lines are circles in pencils of planes
- tangent sphere and tangent cone along curvature lines
- algebraic of degree four, rational of bi-degree $(2,2)$


## Dupin cyclide patches as rational Bézier surfaces



## Supercyclides (E. Blutel, W. Degen)



- projective transforms of Dupin cyclides (essentially)
- conjugate net of conics.
- tangent cones


## Cyclides in CAGD

- surface approximation (Martin, de Pont, Sharrock 1986)
- blending surfaces (Böhm, Degen, Dutta, Pratt, ... ; 1990er)


## Advantages:

- rich geometric structure
- low algebraic degree
- rational parametrization of bi-degree $(2,2)$ :
- curvature line (or conjugate lines)
- circles (or conics)

Dupin cyclides:

- offset surfaces are again Dupin cyclides
- square root parametrization of bisector surface


## Rational parametrization (Dupin cyclides)

Trigonometric parametrization (Forsyth; 1912)

$$
\begin{gathered}
\Phi: f(\theta, \psi)=\frac{1}{a-c \cos \theta \cos \psi}\left(\begin{array}{c}
\mu(c-a \cos \theta \cos \psi)+b^{2} \cos \theta \\
b \sin \theta(a-\mu \cos \psi) \\
b \sin \psi(c \cos \theta-\mu)
\end{array}\right) \\
a, c, \mu \in \mathbb{R} ; b=\sqrt{a^{2}-c^{2}}
\end{gathered}
$$

Representation as Bézier surface

1. $\theta=2 \arctan u, \psi=2 \arctan v$
2. $u \rightsquigarrow \frac{\alpha^{\prime} u+\beta^{\prime}}{\gamma^{\prime} u+\delta^{\prime}}, v \rightsquigarrow \frac{\alpha^{\prime \prime} v+\beta^{\prime \prime}}{\gamma^{\prime \prime} v+\delta^{\prime \prime}}$
3. Conversion to Bernstein basis

## Problem:

A priori knowledge about surface position is necessary (also with other approaches).

## Cyclides as tensor-product Bézier surfaces

Every cyclide patch has a representation as tensor-product Bézier patch of bi-degree ( 2,2 ):
$\mathbf{F}(u, v)=\frac{\sum_{i=0}^{2} \sum_{j=0}^{2} B_{i}^{2}(u) B_{j}^{2}(v) w_{i j} p_{i j}}{\sum_{i=0}^{2} \sum_{j=0}^{2} B_{i}^{2}(u) B_{j}^{2}(v) w_{i j}}, \quad B_{k}^{n}(t)=\binom{n}{k}(1-t)^{n-k} t^{k}$

Aims:

- elementary construction of control points $p_{i j}$
- geometric properties of control net
- elementary construction of weights $w_{i j}$
- applications to CAGD and discrete differential geometry


## The corner points

1. The four corner points $p_{00}, p_{02}, p_{20}$, and $p_{22}$ lie on a circle.


## Reason:

This is true for the prototype parametrizations (torus, circular cone, circular cylinder) and preserved under inversion.

## The missing edge points

2.a The missing edge-points $p_{01}, p_{10}, p_{12}, p_{21}$ lie in the bisector planes of their corner points.
2.b One pair of orthogonal edge tangents can be chosen arbitrarily.

## Reason:

- The edge curves are circles.
- No contradiction because of circularity of edge vertices.

Conclusion
The corner tangent planes envelope a cone of revolution.


## The central control point

3. The central control point $p_{11}$ lies in all four corner tangent planes.

## Reason:

$f(u, v)$ is conjugate parametrization $\Longleftrightarrow$ $f_{u}, f_{v}$ und $f_{u v}$ linear dependent

The quadrilaterals

- $p_{00} p_{01} p_{10} p_{11}$,
- $p_{01} p_{02} p_{12} p_{11}$,
- $p_{10} p_{20} p_{21} p_{10}$,
- $p_{12} p_{21} p_{22} p_{11}$
are planar (conjugate net).



## Parametrization of a circular/conical nets



## Parametrization of a circular/conical nets



## Parametrization of a circular/conical nets



## Parametrization of a circular/conical nets



## Parametrization of a circular/conical nets



## Obvious properties of the control net

Concurrent lines:

- $p_{00} \vee p_{10}$,
- $p_{01} \vee p_{11}$,
- $p_{02} \vee p_{12}$.

Co-axial planes:

- $p_{00} \vee p_{10} \vee p_{20}$,
- $p_{01} \vee p_{11} \vee p_{21}$,
- $p_{02} \vee p_{12} \vee p_{22}$.



## Orthologic tetrahedra

- Non-corresponding sides of the " $x$-axis tetrahedron" and the " $y$-axis tetrahedron" are orthogonal (orthologic tetrahedra).

- The four perpendiculars from the vertices of one tetrahedron on the non-corresponding faces of the other are concurrent.
- Orthology centers are perspective centers for a third tetrahedron.


## The control net as discrete Koenigs-net

- co-planar diagonal points:

$$
\begin{aligned}
& \left(p_{00} \vee p_{11}\right) \cap\left(p_{01} \vee p_{10}\right), \\
& \left(p_{01} \vee p_{12}\right) \cap\left(p_{02} \vee p_{11}\right), \\
& \left(p_{10} \vee p_{21}\right) \cap\left(p_{11} \vee p_{20}\right), \\
& \left(p_{11} \vee p_{22}\right) \cap\left(p_{12} \vee p_{21}\right) .
\end{aligned}
$$

- co-axial planes:

$$
\begin{aligned}
& p_{00} \vee p_{11} \vee p_{02}, \\
& p_{10} \vee p_{11} \vee p_{12}, \\
& p_{20} \vee p_{11} \vee p_{22} .
\end{aligned}
$$

- a net of dual quadrilaterals exists (corresponding edges
 and non-corresponding diagonals are parallel)


Quadrilaterals of vanishing mixed area $\rightsquigarrow$ construction of discrete minimal surfaces.

$$
H=-\frac{A(F, S)}{A(F)}
$$

## The control net of the offset surface



Rich structure comprising circular net, conical net, and three HR congruences:

- existence of offset HR congruence
- existence of orthogonal HR congruence


## The control net of the offset surface



## Literature

击 Degen W.
Generalized cyclides for use in CAGD
In: Bowyer A.D. (editor). The Mathematics of Surfaces IV,
Oxford University Press (1994).
R Huhnen-Venedey E.
Curvature line parametrized surfaces and orthogonal coordinate systems. Discretization with Dupin cyclides Master Thesis, Technische Universität Berlin, 2007.
國 Kaps M.
Teilflächen einer Dupinschen Zyklide in Bézierdarstelllung PhD Thesis, Technische Universität Braunschweig, 1990.

## The weight points

- neighboring control points $p_{i}, p_{j}$
- weights $w_{i}, w_{j}$
- weight point (Farin point)

$$
g_{i j}=\frac{w_{i} p_{i}+w_{j} p_{j}}{w_{i}+w_{j}}
$$



## The weight points

## Properties of weight points

- reconstruction of ratio of weights from weight points is possible
- points in first iteration of rational de Casteljau's algorithm
- weight points of an elementary quadrilateral are necessarily co-planar



## Literature

( Farin, G.
NURBS for Curve and Surface Design - from Projective Geometry to Practical Use
2nd edition, AK Peters, Ltd. (1999)

## Weight points on cyclidic patches



Algorithm of de Casteljau $\Longrightarrow$ weight points of neighboring threads are perspective.

## Weight points on cyclidic patches



Dupin cyclides: One blue and one red weight point can be chosen arbitrarily.

## Weight points on cyclidic patches



Supercyclides: Two blue and two red weight points on neighboring edges can be chosen arbitrarily.

## Determination by edge threads

## Given:

- two edge strips
(control points, weights, apex of tangent cone)
- missing corner point
${ }^{\circ} S_{12}$

$$
\mathrm{s}_{210}
$$

## An auxiliary result

Given are two spatial quadrilaterals with intersecting corresponding edges:

The intersection points points $d_{0}, d_{1}, d_{2}$ und $d_{3}$ are coplanar. $\Longleftrightarrow$
The planes spanned by corre-
 sponding lines intersect in a point.

- The Theorem is self-dual (only one implication needs to be shown).
- If all planes intersect in a point $s$, the two quadrilaterals are perspective with center $s$.


## Dupin cyclide patches

Patch of a Dupin cyclide, bounded by four circular arcs Construction of control points

- Choose four points $p_{00}, p_{02}, p_{22}, p_{20}$ on a circle
- border points $p_{01}, p_{10}, p_{12}, p_{21}$ lie in bisector planes of vertex points
- choose one pair of edge tangents arbitrarily
- find missing border points by reflections
- find central control point as intersection of edge tangent planes


## Open research questions



- (parametrization of asymptotic nets with quadric patches)
- $C^{k}$ conjugate parametrization of conjugate nets
- $C^{k}$ principal parametrization of circular/conical nets and HR-congruences
- parametrization preserving key features of the underlying net

